Exercise Set 3

Exercise 3.1. Consider the following variant of the k-Center problem:

Instance: A complete graph G = (V, E), metric edge weights $d : E(G) \to \mathbb{R}$, a partition $V = C \dot{\cup} S$, an integer $k \in \mathbb{N}$.

Output: A set $X \subseteq S$ with $|X| \le k$ that minimizes

$$\max_{c \in C} \left\{ \min_{s \in X} \{ d(c, s) \} \right\}.$$

- (i) Show that this problem does not admit a $(3-\varepsilon)$ -approximation for any $\varepsilon > 0$ unless P=NP.
- (ii) Give a 3-approximation algorithm.

(4+4 points)

Exercise 3.2. Show that, unless P=NP, for any α polynomial in the input size, there is no α -approximation algorithm for the k-CENTER problem if we do not require the distance function to satisfy the triangle inequality.

(2 points)

Exercise 3.3. Consider the DIRECTED STEINER TREE PROBLEM: Given an edge-weighted digraph G = (V, E), a set of terminals $T \subseteq V$ and a root vertex $r \in V$, find a minimum weight arborescence rooted at r that contains every vertex in T.

Show that a k-approximation algorithm for the DIRECTED STEINER TREE PROBLEM can be used to obtain a k-approximation algorithm for MINIMUM WEIGHT SET COVER.

(2 points)

Exercise 3.4. An instance of MAX-SAT is called k-satisfiable if any k of its clauses can be satisfied simultaneously. Give a polynomial-time algorithm that computes for every 2-satisfiable instance a truth assignment which satisfies at least a $\frac{\sqrt{5}-1}{2}$ -fraction of the clauses.

Hint: Some variables occur in one-element clauses (w.l.o.g. all one-element clauses are positive), set them true with probability a (for some constant $a \in [0,1]$), and set the other variables true with probability $\frac{1}{2}$. Choose a appropriately and derandomize this algorithm.

(4 points)

Deadline: Thursday, April 25^{th} , before the lecture. The websites for lecture and exercises can be found at:

http://www.or.uni-bonn.de/lectures/ss19/appr_ss19_ex.html

In case of any questions feel free to contact me at rockel@or.uni-bonn.de.